

# Entropia

*Até onde uma noção simples pode descrever as  
complexidades?*

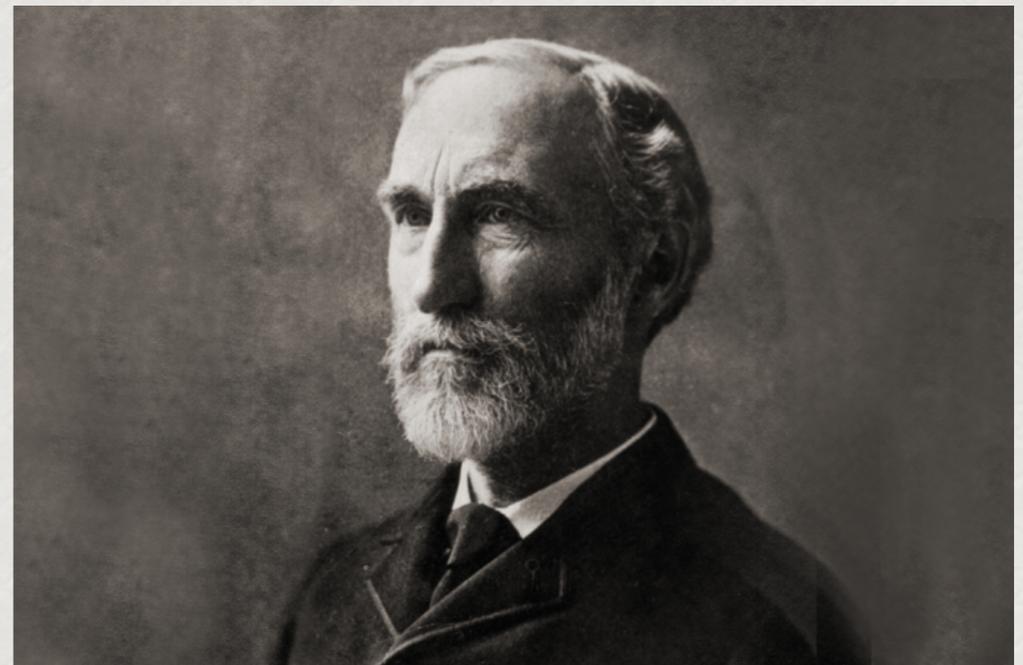
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*Ali Tabzibi*  
*ICMC-USP*

- ◆ *Clausius (1822-88), Ludwig Boltzmann (1844-1906) e J. Willard Gibbs (1839-1903-)*



**L. Boltzmann**



**J. Willard Gibbs:**

# Frases “motivadoras”

*Não há emoção em fenômenos previsíveis!*

*A medida que o tempo passa, conheceremos menos o mundo*

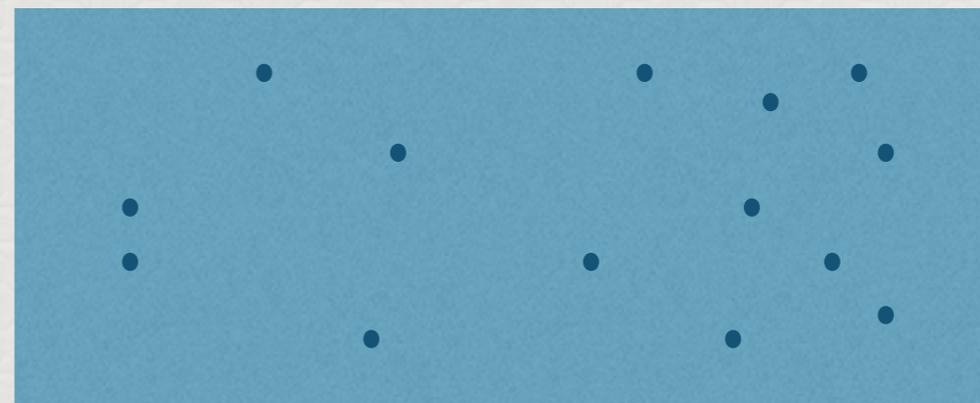
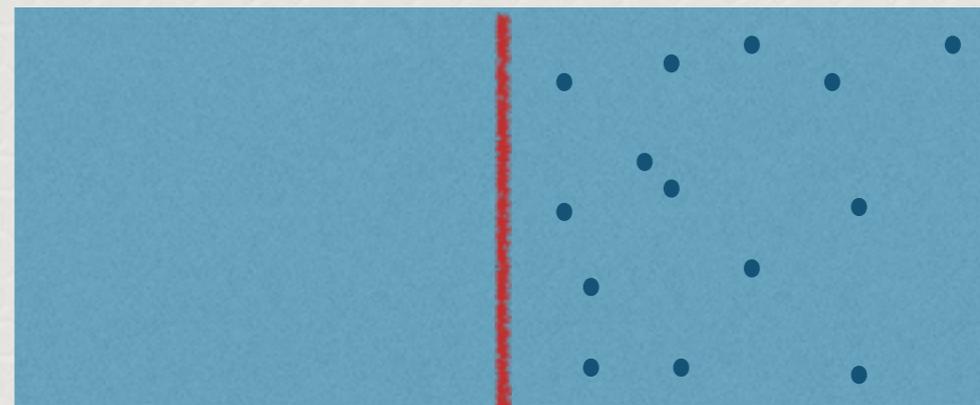
*Porém, não desesperem!*

$$S = k \log (W)$$

*Mecânica Estatística*

*Distribuição de partículas*

$f(t, x, v)$



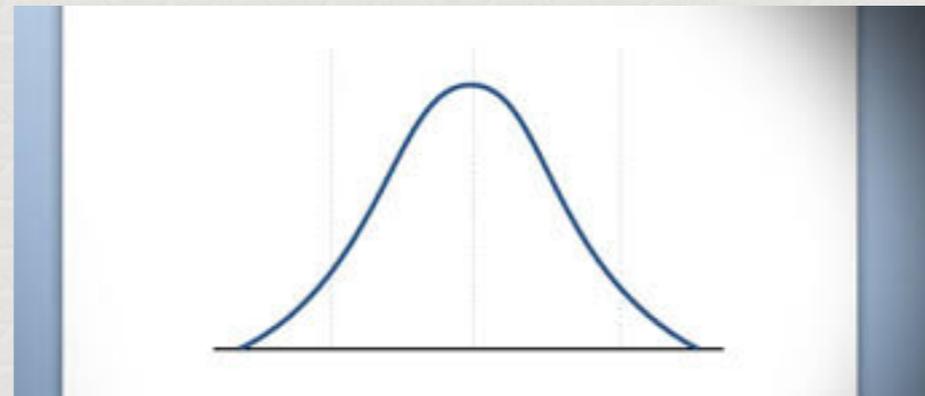
# Boltzmann

$$\frac{\partial f}{\partial t} + v \cdot \nabla_x f = Q(f, f)$$

$$S = -H(f) = - \int f(x, v) \log f(x, v) dv dx$$

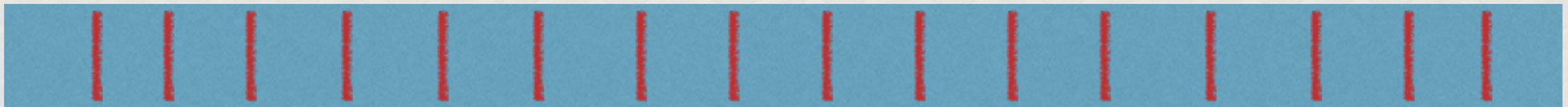
$$S' > 0$$

*Toscani- Villani*



*Stirling*

$$n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$



*K células*

*n: Partículas*

*p<sub>i</sub> Distribuição*

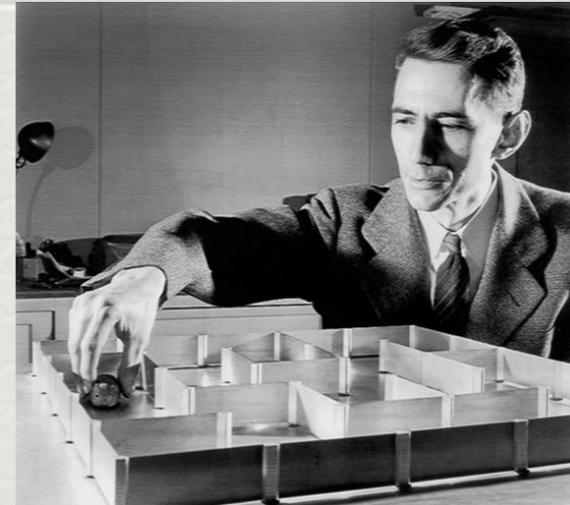
$$\binom{n}{np_1} \binom{n - np_1}{np_2} \cdots \binom{n - np_1 - np_2 - \cdots - np_{k-1}}{np_k}$$

$$= \frac{n!}{(np_1)! \cdots (np_k)!} \sim \frac{1}{p_1^{np_1} p_2^{np_2} \cdots p_k^{np_k}} = e^{nH}$$

$$H = - \sum_{I=1}^k p_i \log(p_i)$$

# Entropia: A incerteza determina valor da informação (Shannon)

*Pai da Teoria de informação*



*Shannon perguntou a Von Neuman, como denotar? ✨*

*Chama de entropia! Ninguém sabe ao certo o que é! Então ✨  
nas discussões ganhará!*

# Shannon entropy

$$I(\Lambda) = -k \log(\mu(\Lambda))$$

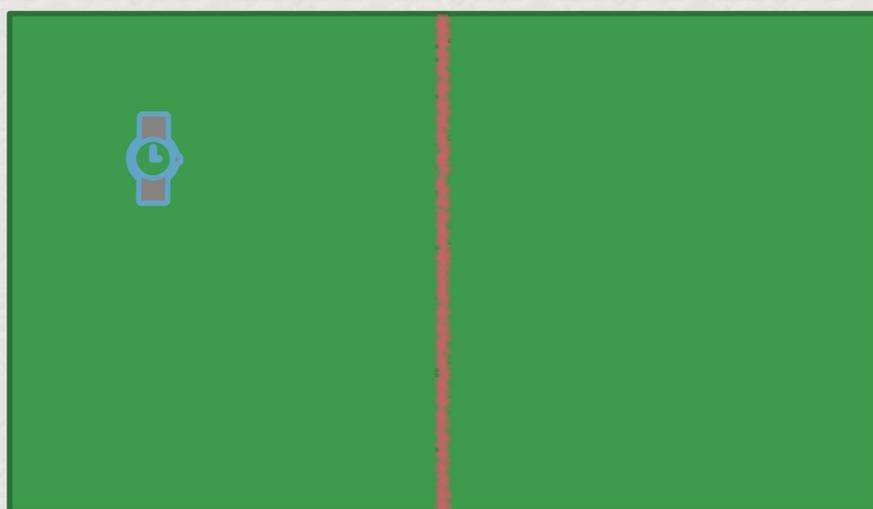
*Um fenômeno com probabilidade total tem incerteza zero.* ✨

*Pouca probabilidade, mais incerteza* ✨

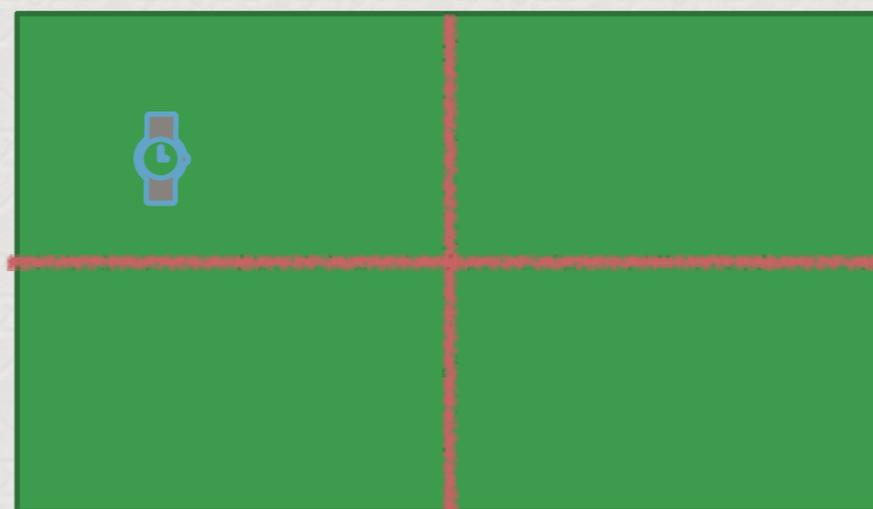
*Fenômenos independentes, somam incertezas* ✨

# *Chave perdida no campo*

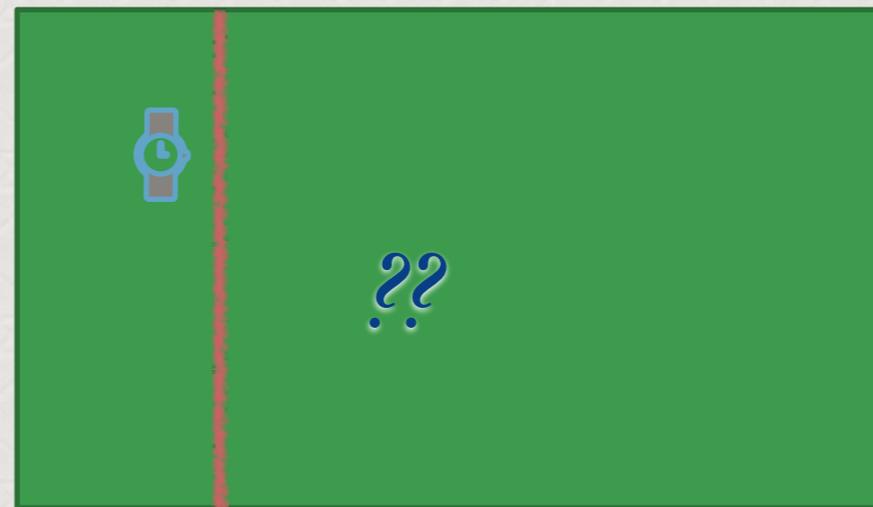
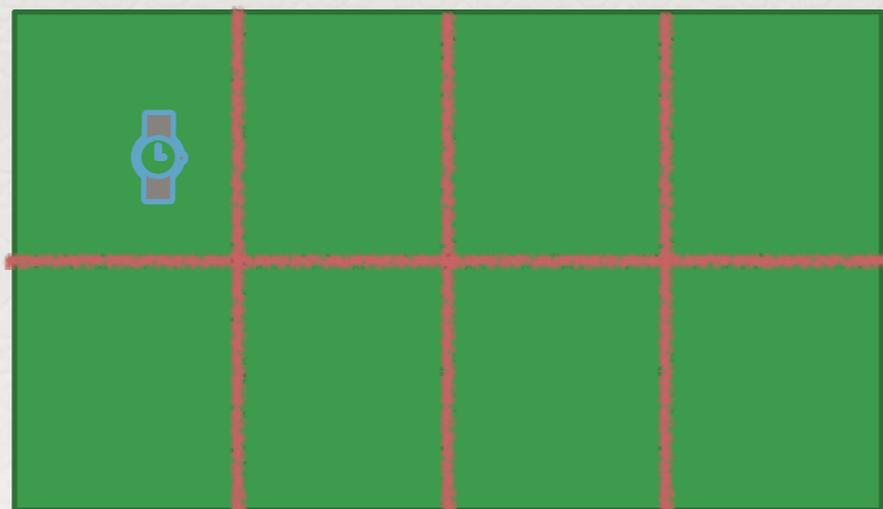
*1*



*2*



*3*



*0.8I...*

# Shannon entropy



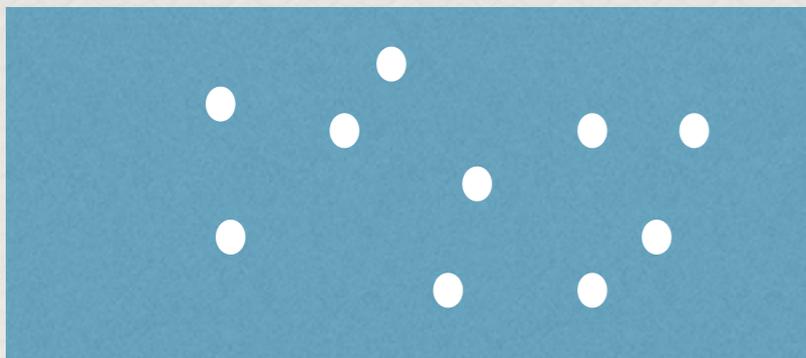
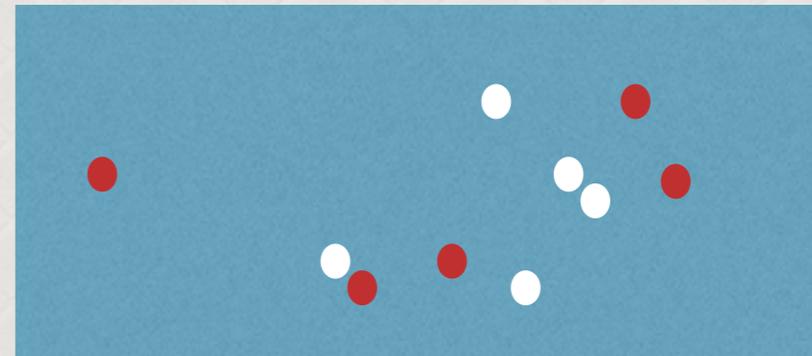
$$\mathcal{P} = \{P_1, P_2, \dots, P_n\}$$

$$H(\mathcal{P}) = - \sum p_i \log(p_i)$$

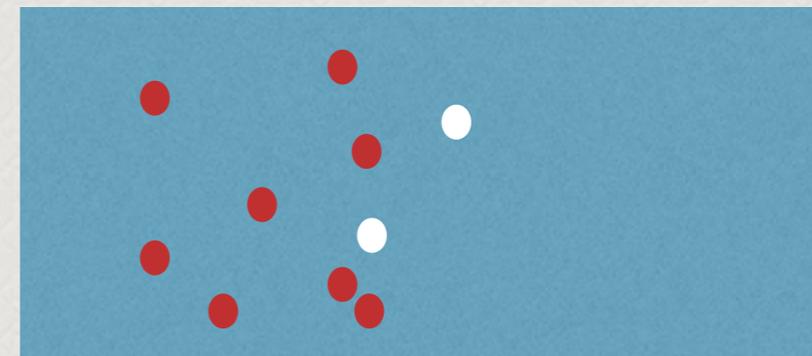
# *Informação escondida numa partição*

## *A cor sobre número da bola!*

$$-(1/2 \log(1/2) + 1/2 \log(1/2)) = 1$$

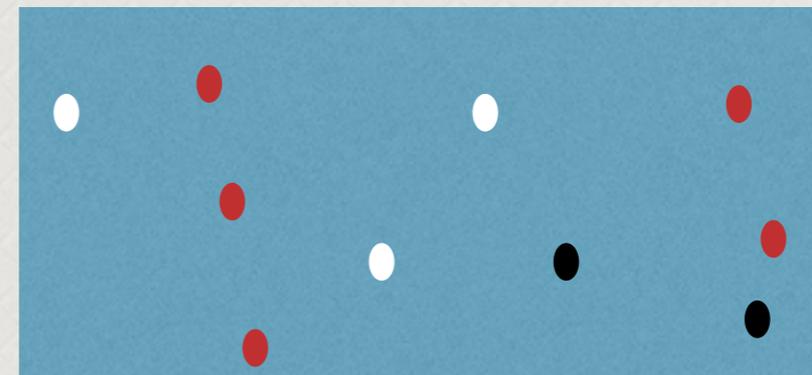


< 1



0

$$-(0.2 \log(0.2) + 0.3 \log(0.3) + 0.5 \log(0.5)) \sim 1.49$$



# Multiplicadores de Lagrange

$$f(p_1, \dots, p_n) = \sum_{i=1}^n p_i \log p_i$$

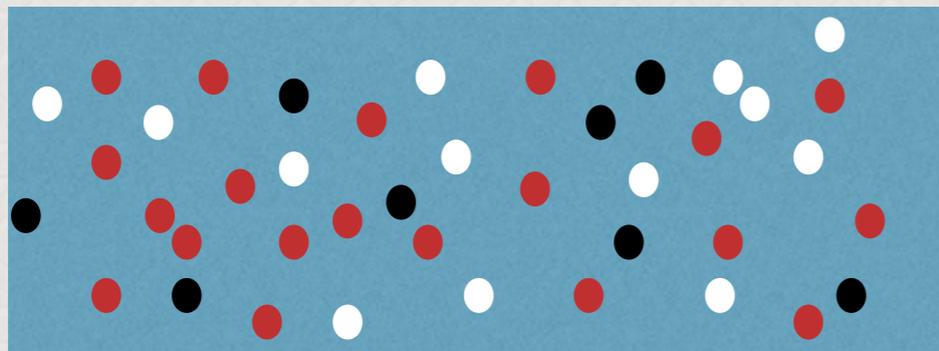
$$g(p_1, \dots, p_n) = \sum_{i=1}^n p_i$$

$$Df = -(1 + \log(p_1), \dots, 1 + \log(p_n))$$

$$Df = \lambda Dg$$

$$p_i = p_j \Rightarrow p_i = 1/n.$$

# Jogo de Sim ou Não



●	0.5	<i>Prob.</i>
●	0.2	
●	0.3	

انتخاب یک توپ به طور تصادفی و کشف رنگ توپ با پرسشهای از نوع بله و خیر

*P1: A bola é preta ou branca?*

*Sim: P2*  
*Não: Vermelha*

*P2: A bola é preta?*

*Dividir o conjunto em dois subconjuntos com a probabilidade mais próximo possível  $P_1, P_2$*

*Dividir a resposta da  $P_1$  em dois subconjuntos com probabilidade mais próximo possível*

*Continuar o algoritmo até chegar apenas duas cores. ...*

# Número de perguntas e entropia

$$-(0.2\log(0.2) + 0.3\log(0.3) + 0.5\log(0.5)) \sim 1.49 \quad 0.5 + 0.2 \times 2 + 0.3 \times 2 = 1.5$$

$N_i$ : Número de perguntas até chegar a cor  $i$

$$h = - \sum p_i \log_2 p_i \leq \sum p_i N_i$$

*Entropia  $\leq$  Média de perguntas*

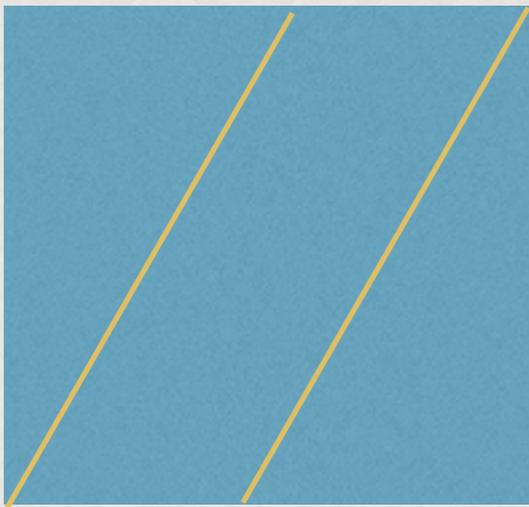
# Sistemas Dinâmicos



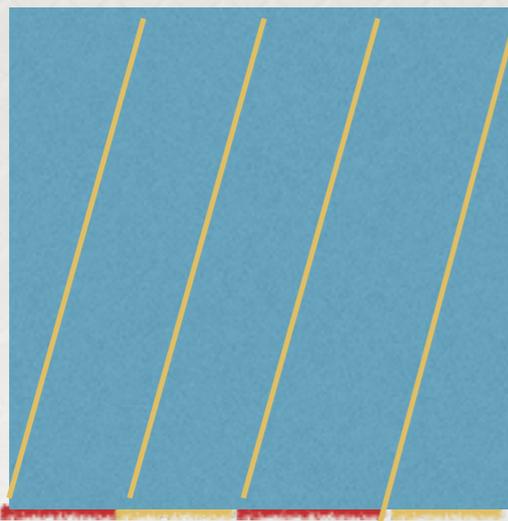
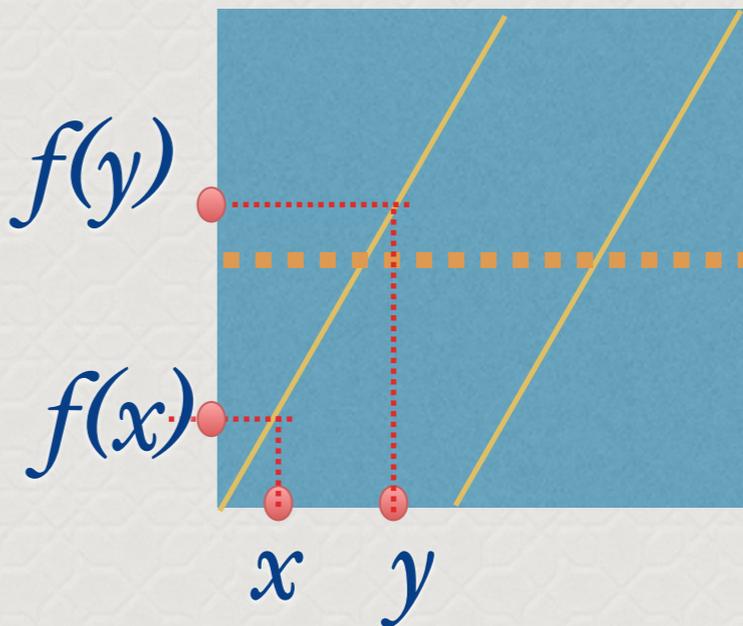
- ✦ *Kolmogorov (1903- 1987), (Kakutani), Sinai (1935 , -)*
- ✦ *Sinai (Prêmio Abel 2014).*

*Primeiramente pensava que entropia apenas separa estocástico de determinístico!*

# Sistemas Dinâmicos



$$f(x) = 2x(\text{mod } -1)$$



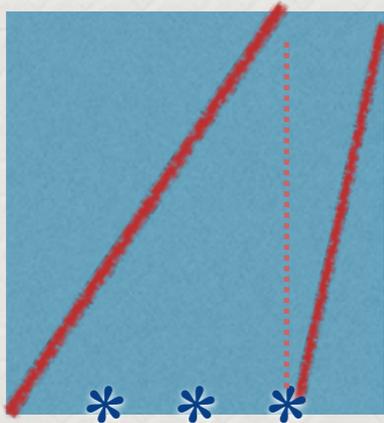
*Taxa de criação de informação*

$$\log(4) - \log(2) = \log(2)$$

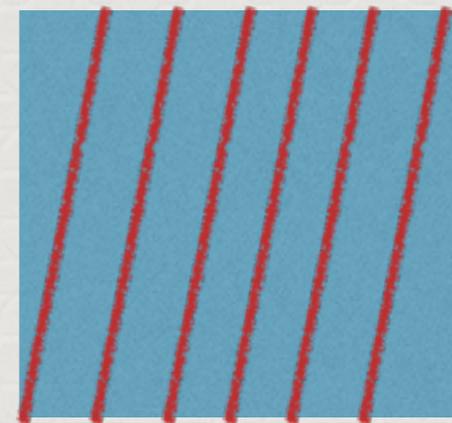
$$\log(2^n) - \log(2^{n-1}) = \log(2)$$

# Lançamento de moeda ou Roda da sorte

$$h = -\sum p_i \log(p_i)$$

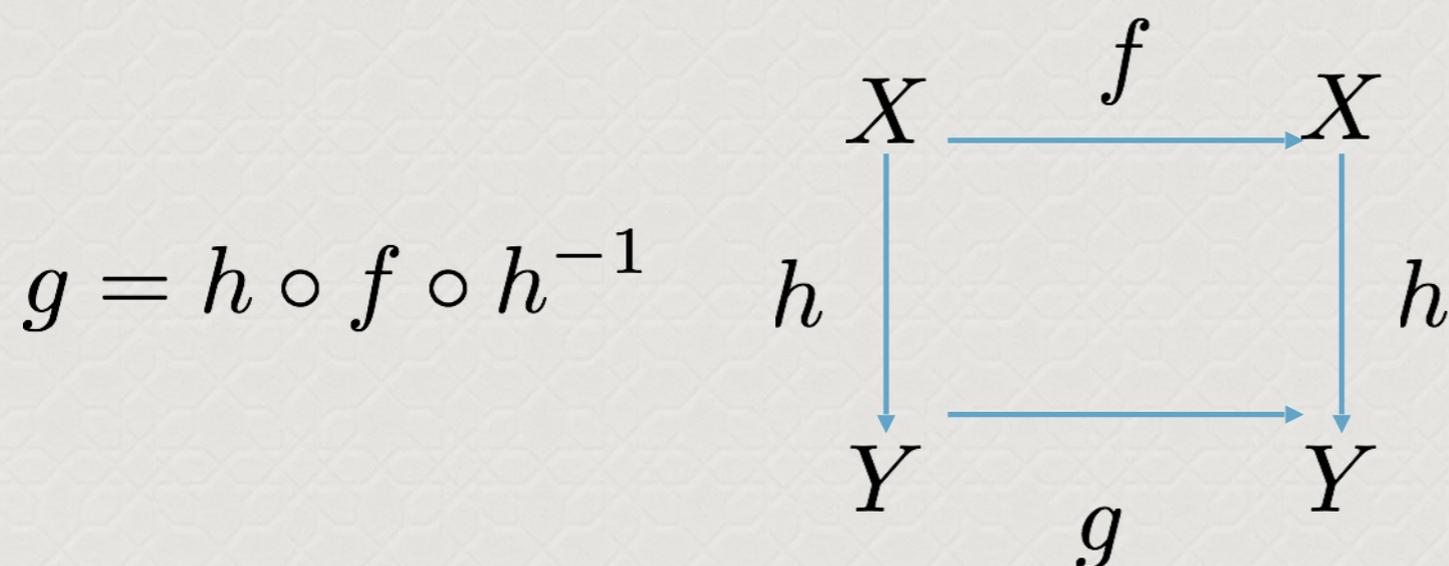


$0.81$



$\log_2(6) \sim 2.58$

# Conjugação Dinâmica



*Topologicamente conjugado : homeomorfismo  $h$*

*Metricamente conjugado: Preservando medida  $h$*

# Entropia como invariante

*Sistemas Conjugados tem entropias iguais.*

*Não é invariante completa:*

*Identidade* ✪

*Rotação* ✪

# Ornstein



- *Ornstein: Duas rodas da sorte com mesma entropia são conjugadas metricamente.*
- *Entropia é invariante completo no mundo dos sistemas mais aleatório possível.*

## *Entropia métrica*

$$B_n(x, \epsilon) := \{y \mid d(f^n(x), f^n(y)) \leq \epsilon\}$$

$$\lim_{\epsilon \rightarrow 0} \limsup_{n \rightarrow \infty} \frac{-1}{n} \log \mu(B_n(x, \epsilon))$$

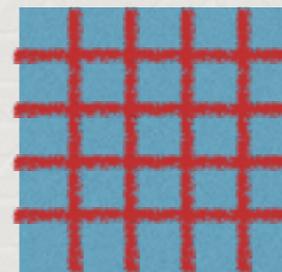
*Exemplo*  $2x \pmod{-1}$       $\mu(B_n(x, \epsilon)) = \frac{2\epsilon}{2^n}$

# Dimensão

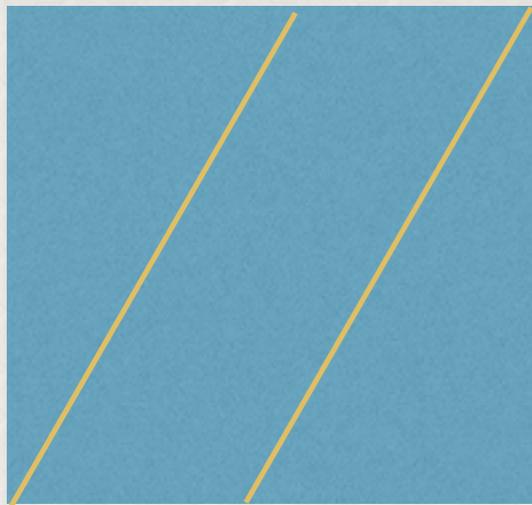
*Dimensão de um conjunto é o “número” das informações necessárias para localizar os pontos do conjunto.* ✪

*Capacity dimension:*  $\dim_K(S)$

$$- \limsup_{\epsilon \rightarrow 0} \frac{\log(N(\epsilon))}{\log \epsilon}$$



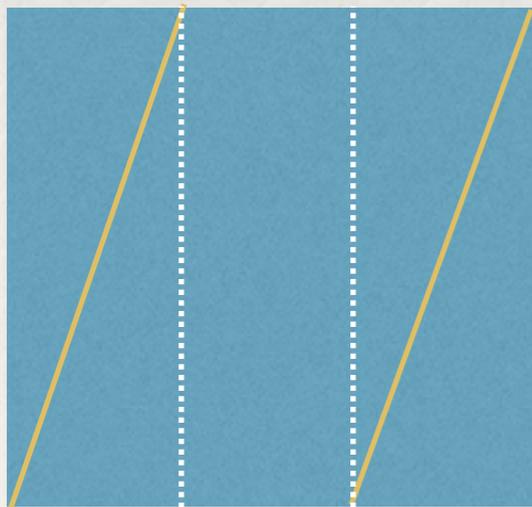
*Cobertura ótima com abertos de diâmetro menor que  $\epsilon$*  €



$$f(x) = 2x(\text{mod } 1)$$

$\log(2)$  **Expoente de expansão**

$\log(2)$  *entropia*



$1/3$   $2/3$

*Conjunto de Cantor*

$\log(3)$  **Expansão**

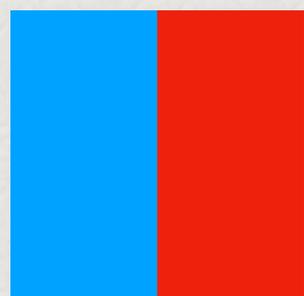
$\log(2)$  *Entropia*

$\log(2)/\log(3)$  *Dimensão*

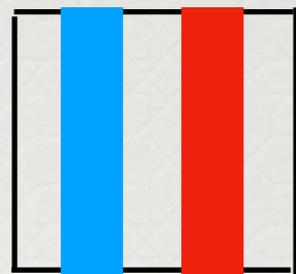
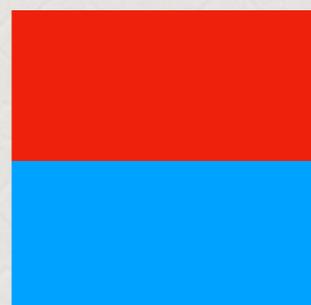
# Ruelle

$$h_\mu \leq \sum_{\lambda_i > 0} \lambda_i$$

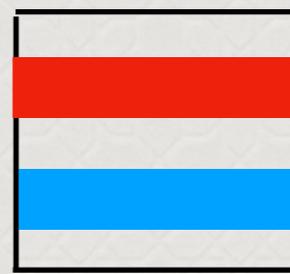
$\mu$



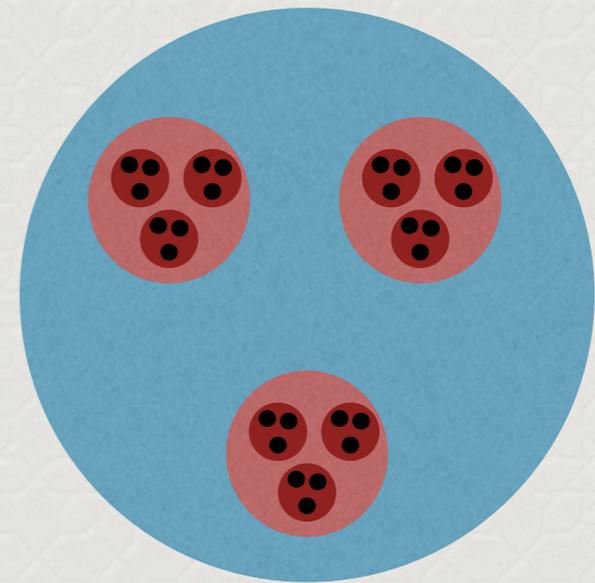
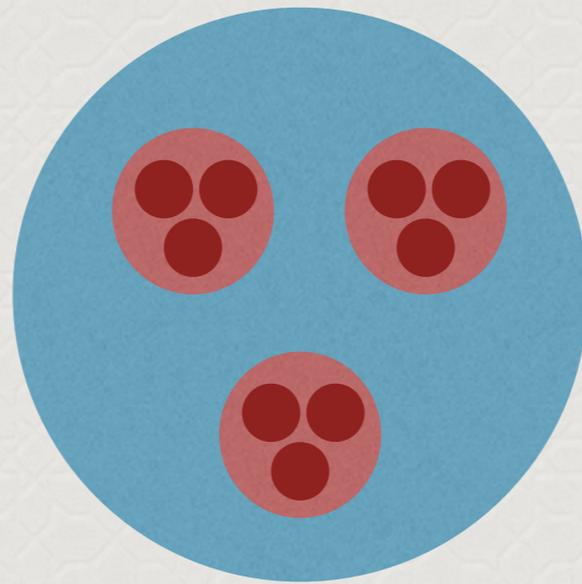
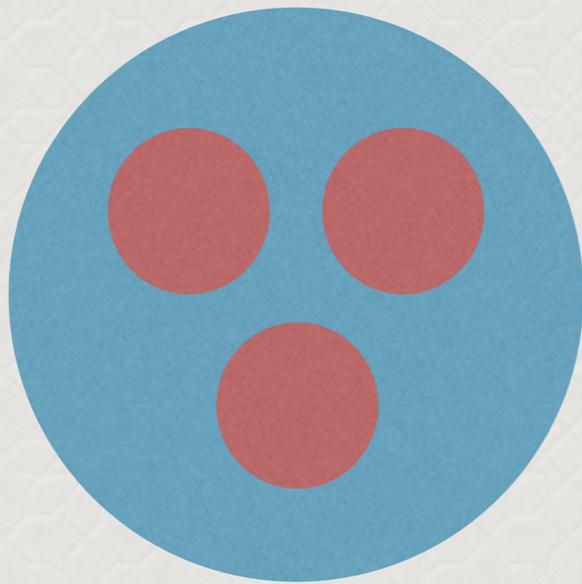
$$\xrightarrow{\mathbb{F}} h_\mu = \lambda = \log(2)$$



$$\xrightarrow{\mathcal{G}} h_\mu = \log(2) < \lambda$$



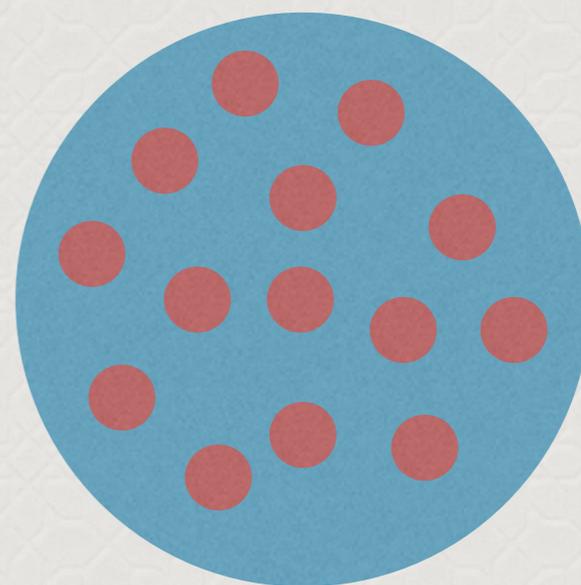
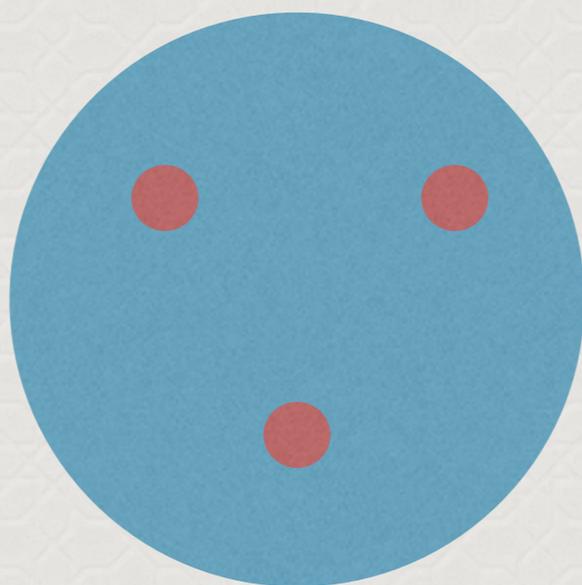
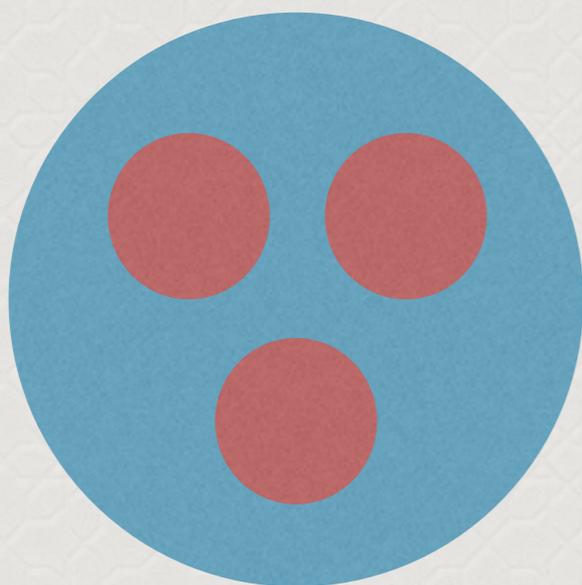
# Dimensão $\times$ Expansão = Entropia



$$F : \bigcup_{i=1}^3 B_i \rightarrow B$$

$$\bigcap_{n \geq 0} f^{-n}(B)$$

انبساط × بعد = أنتروپی



*Obrigado pela  
atenção*

